

# New CAP Reduction Mechanisms for IEEE 802.15.4 DSME to Support Fluctuating Traffic in IoT Systems

Florian Meyer and Ivonne Mantilla-Gonzalez and Volker Turau

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# Agenda

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- Motivation

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- Description of the New CAP reduction Mechanisms

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- Theoretical and simulative evaluation

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- Conclusions and Outlook

# Motivation

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- Efficient management of resources needed to avoid
  - ◆ Channel underutilization
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- IEEE 802.15.4 DSME  $\Rightarrow$  reliability, scalability and energy efficiency in IoT applications
  - ◆ TDMA/FDMA reservation scheme
  - ◆ Distributed slot negotiation and collision free slot allocation
  - ◆ CAP reduction mode
  - ◆ Fixed trade-off between adaptability and throughput

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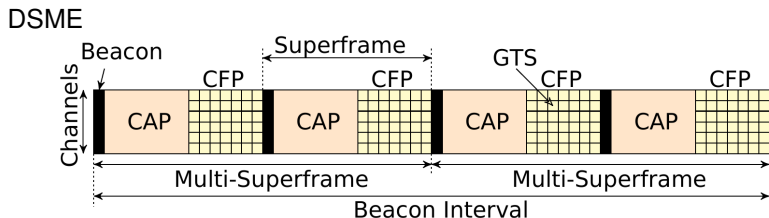
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$\Rightarrow$  **We propose two extensions of DSME to provide high degree of responsiveness to traffic fluctuations while keeping throughput high**

# Frame Structure

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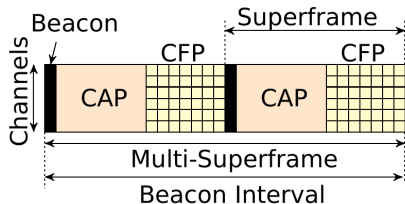


## Parameters:

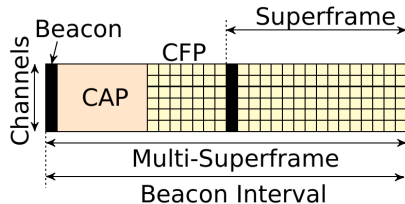
- Contention Access Period : CAP
- Contention Free Period : CFP
- Guaranteed Time slot : GTS
- Superframe order ( $SO$ ):  
Length of a slot / superframe
- Multi-Superframe order ( $MO$ ):  
Length of a Multi-Superframe
- Beacon order ( $BO$ ):  
Length of a beacon interval (BI)

# Frame Structure

DSME - no CAP reduction (NCR)

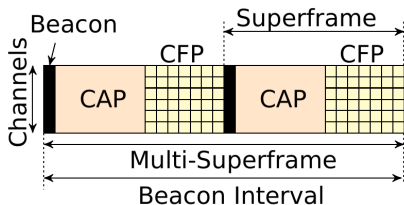


DSME - CAP reduction (CR)

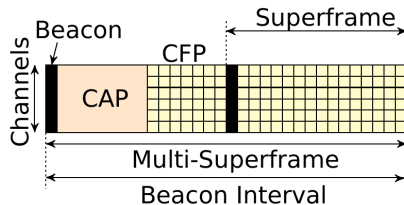


# Frame Structure

DSME - no CAP reduction (NCR)



DSME - CAP reduction (CR)



**Definition: Fraction  $\tau$  of CFPs time slots in a beacon interval**

$$\tau = \frac{\text{Total number of CFP's time slots in a BI}}{\text{Total number of time slots in a BI}}$$

In this case,  $\tau_{\text{NCR}} = 7/16$  and  $\tau_{\text{CR}} = 11/16$

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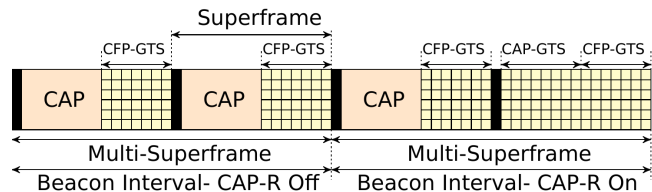
In DSME  $\tau$  is set before deployment and cannot be changed at run-time

## *Challenge:*

How to extend DSME such that  $\tau$  can be changed dynamically with a fine granularity?

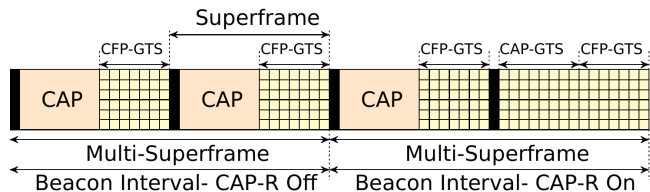
# Alternating CAP reduction - ACR

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Alternation between CR and NCR every beacon interval

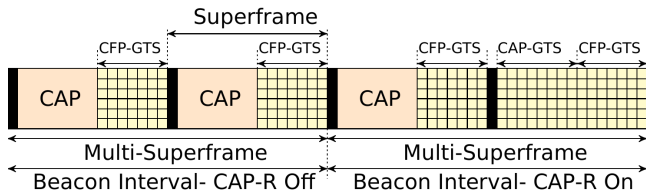
# Alternating CAP reduction - ACR



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**How is this achieved?**

# Alternating CAP reduction - ACR



Alternation between CR and NCR every beacon interval

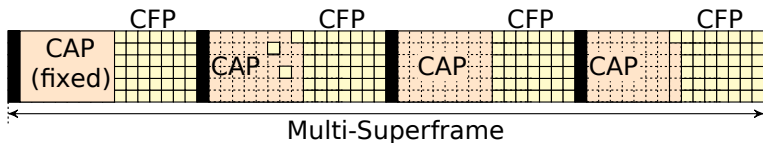
## How is this achieved?

- PAN coordinator defines current operating mode per BI
- Operating mode is encoded in beacon messages
- Synchronization in network is guaranteed with minimum effect on CSMA traffic

# Dynamic CAP reduction - DCR

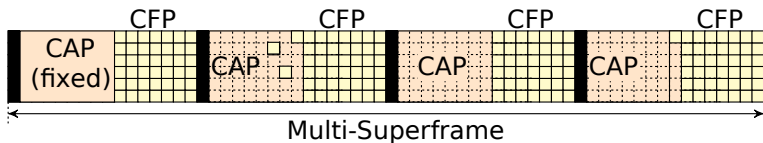


# Dynamic CAP reduction - DCR



Allocation of GTTs in CAPs locally according to node's traffic demands

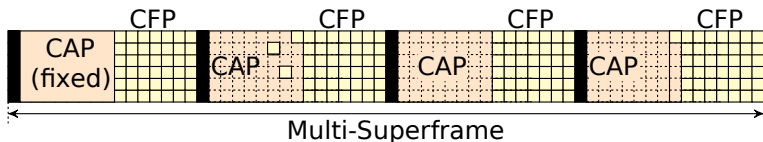
# Dynamic CAP reduction - DCR



Allocation of GTTs in CAPs locally according to node's traffic demands

**How is this achieved?**

# Dynamic CAP reduction - DCR



Allocation of GTSSs in CAPs locally according to node's traffic demands

## How is this achieved?

- DCR starts in NCR and dynamically shrinks CAPs until CR mode
- Shrinking mechanism begins once all GTSSs in CFPs are used
- Channel for transmissions in reduced CAPs  $\neq Channel_{CAP}$

# Theoretical evaluation

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Values of  $\tau$  for different values of  $MO$ ,  $BO = 7$  and  $SO = 3$

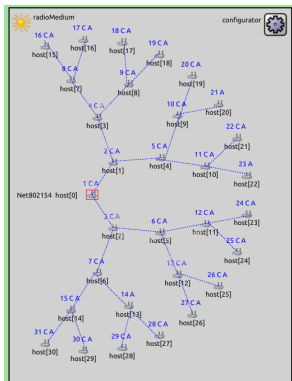
	$MO = 4$	$MO = 5$	$MO = 6$	$MO = 7$
NCR	43.75%	43.75%	43.75%	43.75%
CR	68.75%	81.25%	87.5%	90.06%
ACR	56.25%	62.5%	65.63%	67.19%
DCR	43.75% - 68.75%	43.75% - 81.25%	43.75% - 87.5%	43.75% - 90.06%

# Theoretical evaluation

Expected time to send a CAP message on time slot level for different values of  $MO$ ,  $BO = 7$  and  $SO = 3$

	$MO = 4$	$MO = 5$	$MO = 6$	$MO = 7$
NCR	2.25	2.25	2.25	2.25
CR	9.38	24.94	56.72	120.61
ACR	5.81	14.10	29.49	61.43
DCR	2.25 - 9.38	2.25 - 24.94	2.25 - 56.72	2.25 - 120.61

# Simulative evaluation



- Data collection - convergecast pattern
- Rooted binary tree (31 nodes)
- Packet generation rate  $\Rightarrow \delta$  [packets/s]
- Packet generation  $\Rightarrow$  Poisson distribution with mean =  $\lambda$ 
  - ◆ Varying burst sizes
  - ◆ Varying  $\delta$

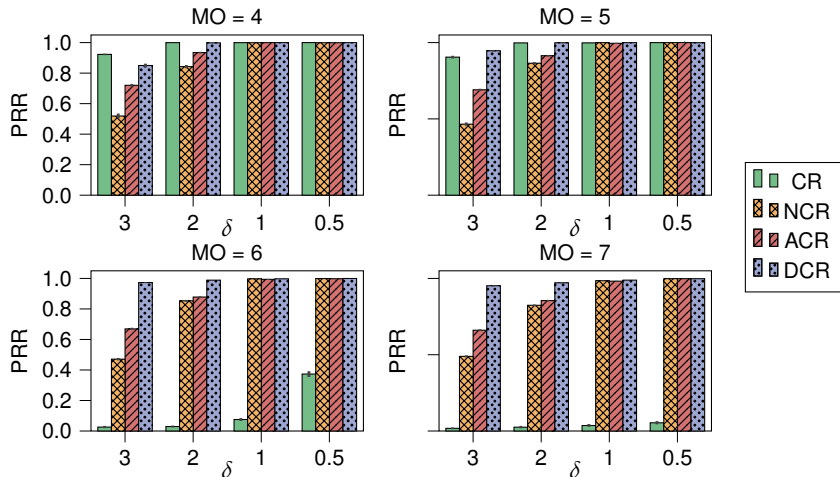


Parameter	SO	MO	BO	$Q_{GTS}$	$\delta$
Values	3	{4,...,7}	7	22	{1,...,4}

# Simulative evaluation

Scenario: varying  $\delta$

Packet Reception Ratio (PRR)





# Conclusions and Outlook

## *Challenge:*

How to extend DSME such that  $\tau$  can be dynamically adjusted after deployment of the network?

- **ACR**  $\Rightarrow$  combines strengths of CR and NCR. It is far easier to implement and remains within the original standard
- **DCR**  $\Rightarrow$  is more flexible with respect to  $\tau$ . It does not conform with the IEEE 802.15.4 standard

## **Future work**

- To fully exploit the possibilities of ACR and DCR  $\Rightarrow$  a powerful dynamic scheduler is required

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